The Karnaugh Map

COE 202

Digital Logic Design

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Presentation Outline

- Boolean Function Minimization
- The Karnaugh Map (K-Map)
- Two, Three, and Four-Variable K-Maps
- Prime and Essential Prime Implicants
- Minimal Sum-of-Products and Product-of-Sums
- Don't Cares
- Five and Six-Variable K-Maps
- Multiple Outputs

Boolean Function Minimization

- Complexity of a Boolean function is directly related to the complexity of the algebraic expression
- The truth table of a function is unique
- However, the algebraic expression is not unique
- Boolean function can be simplified by algebraic manipulation
- However, algebraic manipulation depends on experience
- Algebraic manipulation does not guarantee that the simplified Boolean expression is minimal

Example: Sum of Minterms

Truth Table

x y z	f	Minterm
0 0 0	0	
0 0 1	1	$m_1 = x'y'z$
0 1 0	1	$m_2 = x'yz'$
0 1 1	1	$m_3 = x'yz$
1 0 0	0	
1 0 1	1	$m_5 = xy'z$
1 1 0	0	
1 1 1	1	$m_7 = xyz$

Focus on the '**1**' entries

$$f = m_1 + m_2 + m_3 + m_5 + m_7$$

$$f = \sum(1, 2, 3, 5, 7)$$

$$f = x'y'z + x'yz' + x'yz + xy'z + xyz$$

- Sum-of-Minterms has 15 literals ➔ Can be simplified

Algebraic Manipulation

- **Simplify:** $f = x'y'z + x'yz' + x'yz + xy'z + xyz$ (15 literals)

$f = x'y'z + x'yz' + x'yz + xy'z + xyz$	(Sum-of-Minterms)
$f = x'y'z + x'yz + x'yz' + xy'z + xyz$	Reorder
$f = x'z(y' + y) + x'yz' + xz(y' + y)$	Distributive $\cdot$ over +
$f = x'z + x'yz' + xz$	Simplify (7 literals)
$f = x'z + xz + x'yz'$	Reorder
$f = (x' + x)z + x'yz'$	Distributive $\cdot$ over +
$f = z + x'yz'$	Simplify (4 literals)
$f = (z + x'y)(z + z')$	Distributive + over $\cdot$
$f = z + x'y$	Simplify (3 literals)

Drawback of Algebraic Manipulation

- No clear steps in the manipulation process
  - Not clear which terms should be grouped together
  - Not clear which property of Boolean algebra should be used next
- Does not always guarantee a minimal expression
  - Simplified expression may or may not be minimal
  - Different steps might lead to different non-minimal expressions
- However, the goal is to minimize a Boolean function
- Minimize the **number of literals** in the Boolean expression
  - The **literal count** is a good measure of the **cost** of logic implementation
  - Proportional to the number of transistors in the circuit implementation

Karnaugh Map

- ❖ Called also K-map for short
- ❖ The Karnaugh map is a diagram made up of squares
- ❖ It is a reorganized version of the truth table
- ❖ Each square in the Karnaugh map represents a minterm
- ❖ Adjacent squares differ in the value of one variable
- ❖ Simplified expressions can be derived from the Karnaugh map
  - ✧ By recognizing patterns of squares
- ❖ Simplified sum-of-products expression (AND-OR circuits)
- ❖ Simplified product-of-sums expression (OR-AND circuits)

Next . . .

- ❖ Boolean Function Minimization
- ❖ The Karnaugh Map (K-Map)
- ❖ Two, Three, and Four-Variable K-Maps
- ❖ Prime and Essential Prime Implicants
- ❖ Minimal Sum-of-Products and Product-of-Sums
- ❖ Don't Cares
- ❖ Five and Six-Variable K-Maps
- ❖ Multiple Outputs

Two-Variable Karnaugh Map

- Minterms m_0 and m_1 are adjacent (also, m_2 and m_3)
  - They differ in the value of variable y
- Minterms m_0 and m_2 are adjacent (also, m_1 and m_3)
  - They differ in the value of variable x

Two-variable K-map

$x \backslash y$	0	1
0	m_0	m_1
1	m_2	m_3

$x \backslash y$	0	1
0	$x'y'$	$x'y$
1	xy'	xy

From a Truth Table to Karnaugh Map

- Given a truth table, construct the corresponding K-map
- Copy the function values from the truth table into the K-map
- Make sure to copy each value into the proper K-map square

Truth Table

x y	f
0 0	1
0 1	0
1 0	1
1 1	1

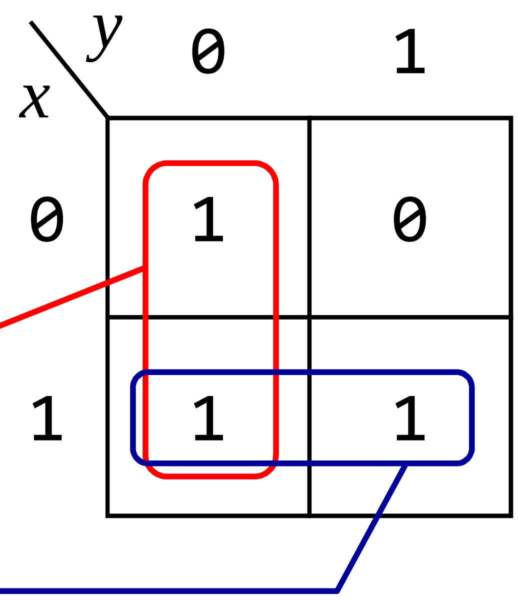
K-map

$x \backslash y$	0	1
0	1	0
1	1	1

K-Map Function Minimization

- ❖ Two adjacent cells containing 1's can be combined
- ❖ $f = m_0 + m_2 + m_3$
- ❖ $f = x'y' + xy' + xy$ (6 literals)
- ❖ $m_0 + m_2 = x'y' + xy' = (x' + x)y' = y'$
- ❖ $m_2 + m_3 = xy' + xy = x(y' + y) = x$
- ❖ Therefore, f can be simplified as: $f = x + y'$ (2 literals)

K-map



Three-Variable Karnaugh Map

- Have eight squares (for the 8 minterms), numbered 0 to 7
- The last two columns are not in numeric order: 11, 10
  - Remember the numbering of the squares in the K-map
- Each square is adjacent to three other squares
- Minterms in adjacent squares can always be combined
  - This is the key idea that makes the K-map work
- Labeling of rows and columns is also useful

$x \backslash yz$	00	01	11	10
0	m_0	m_1	m_3	m_2
1	m_4	m_5	m_7	m_6

	$x \backslash yz$	00	01	11	10
		y'	y'	y	y
x'	0	$x'y'z'$	$x'y'z$	$x'yz$	$x'yz'$
x	1	$xy'z'$	$xy'z$	xyz	xyz'
		z'	z	z	z'

Simplifying a Three-Variable Function

Simplify the Boolean function: $f(x, y, z) = \sum(3, 4, 5, 7)$

$f = x'yz + xy'z' + xy'z + xyz$ (12 literals)

1. Mark **'1'** all the K-map squares that represent function f

2. Find possible adjacent squares

x \ yz	00 (y')	01 (y')	11 (y)	10 (y)
x' 0	0	0	1	0
x 1	1	1	1	0
	z'	z	z	z'

$x'yz + xyz = (x' + x)yz = yz$

$xy'z' + xy'z = xy'(z' + z) = xy'$

Therefore, $f = xy' + yz$ (4 literals)

Simplifying a Three-Variable Function (2)

Here is a second example: $f(x, y, z) = \sum(3, 4, 6, 7)$

$f = x'yz + xy'z' + xyz' + xyz$ (12 literals)

Learn the locations of the 8 indices based on the variable order

$x'yz + xyz = (x' + x)yz = yz$

Corner squares can be combined

$xy'z' + xyz' = xz'(y' + y) = xz'$

Therefore, $f = xz' + yz$ (4 literals)

		y'		y	
x \ yz		00	01	11	10
x'	0	0	0	1	0
x	1	1	0	1	1
		z'	z	z	z'

Combining Squares on a 3-Variable K-Map

- By combining squares, we reduce number of literals in a product term, thereby reducing the cost
- On a 3-variable K-Map:
  - One square represents a minterm with 3 variables
  - Two adjacent squares represent a term with 2 variables
  - Four adjacent squares represent a term with 1 variable
  - Eight adjacent square is the constant **'1'** (no variables)

Example of Combining Squares

- Consider the Boolean function: $f(x, y, z) = \sum(2, 3, 5, 6, 7)$
- $f = x'yz' + x'yz + xy'z + xyz' + xyz$
- The four minterms that form the 2×2 red square are reduced to the term y
- The two minterms that form the blue rectangle are reduced to the term xz
- Therefore: $f = y + xz$

		y'	y'	y	y
$x \backslash yz$		00	01	11	10
x'	0	0	0	1	1
x	1	0	1	1	1
		z'	z	z	z'

$$x'yz + x'yz' + xyz + xyz'$$
$$= x'y(z + z') + xy(z + z')$$
$$= x'y + xy = (x' + x)y = y$$

Minimal Sum-of-Products Expression

Consider the function: $f(x, y, z) = \sum(0, 1, 2, 4, 6, 7)$

Find a minimal sum-of-products (SOP) expression

Solution:

Red block: term = z'

Green block: term = $x'y'$

Blue block: term = xy

$x \backslash yz$	00 (y', z')	01 (y', z)	11 (y, z)	10 (y, z')
x' 0	1	1	0	1
x 1	1	0	1	1

Minimal sum-of-products: $f = z' + x'y' + xy$ (5 literals)

Four-Variable Karnaugh Map

4 variables ➔ 16 squares

Remember the numbering of the squares in the K-map

Each square is adjacent to four other squares

$$
\begin{aligned}
m_0 &= w'x'y'z' & \quad m_1 &= w'x'y'z \\
m_2 &= w'x'y\,z' & \quad m_3 &= w'x'y\,z \\
m_4 &= w'x\,y'z' & \quad m_5 &= w'x\,y'z \\
m_6 &= w'x\,y\,z' & \quad m_7 &= w'x\,y\,z \\
m_8 &= w\,x'y'z' & \quad m_9 &= w\,x'y'z \\
m_{10} &= w\,x'yz' & \quad m_{11} &= w\,x'y\,z \\
m_{12} &= w\,x\,y'z' & \quad m_{13} &= w\,x\,y'z \\
m_{14} &= w\,x\,y\,z' & \quad m_{15} &= w\,x\,y\,z
\end{aligned}
$$

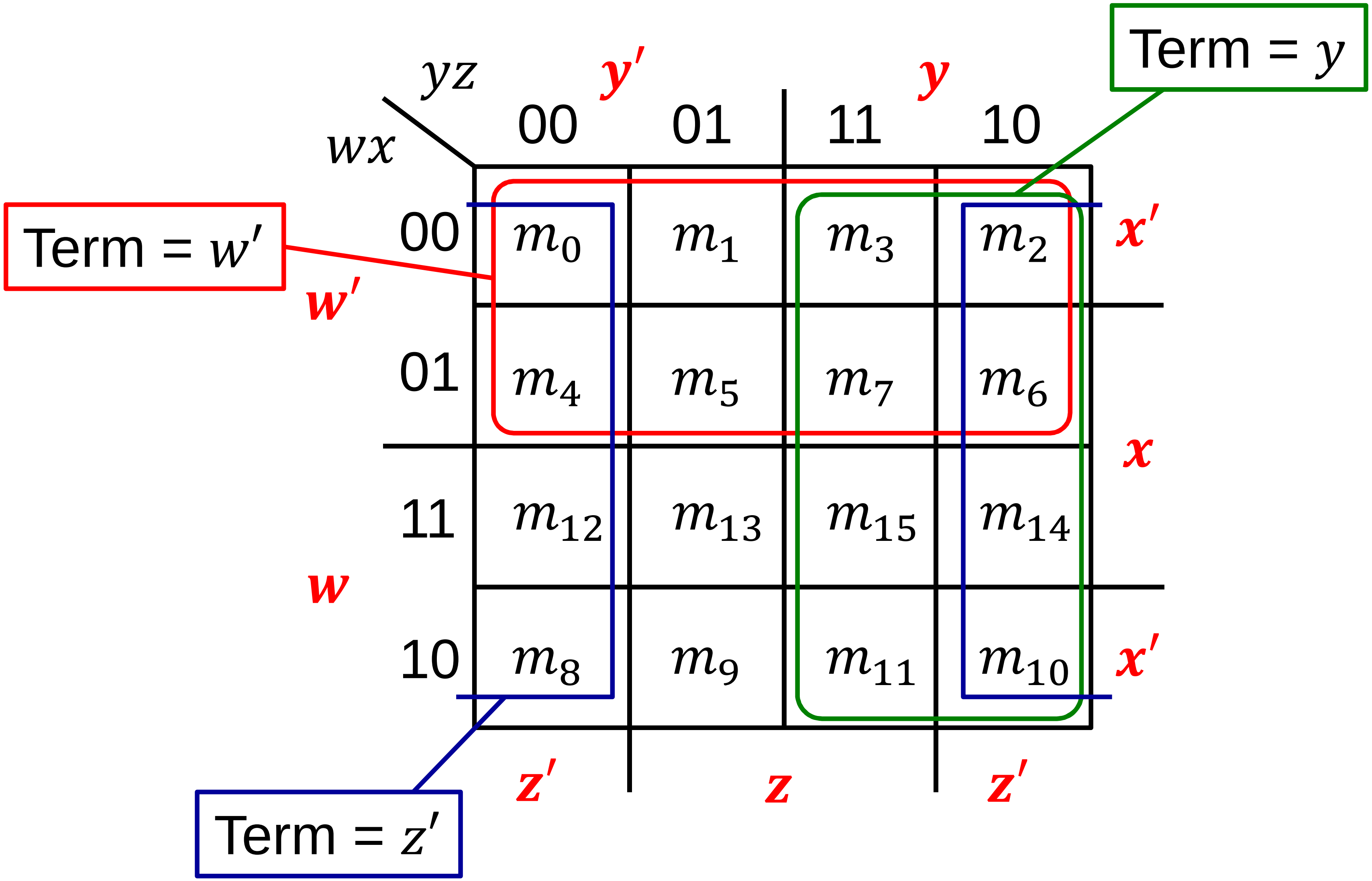
Notice the order of Rows 11 and 10 and the order of columns 11 and 10

	$wx \backslash yz$	00	01	11	10	
		y'		y		
w'	00	m_0	m_1	m_3	m_2	x'
	01	m_4	m_5	m_7	m_6	
w	11	m_{12}	m_{13}	m_{15}	m_{14}	x
	10	m_8	m_9	m_{11}	m_{10}	x'
		z'	z		z'	

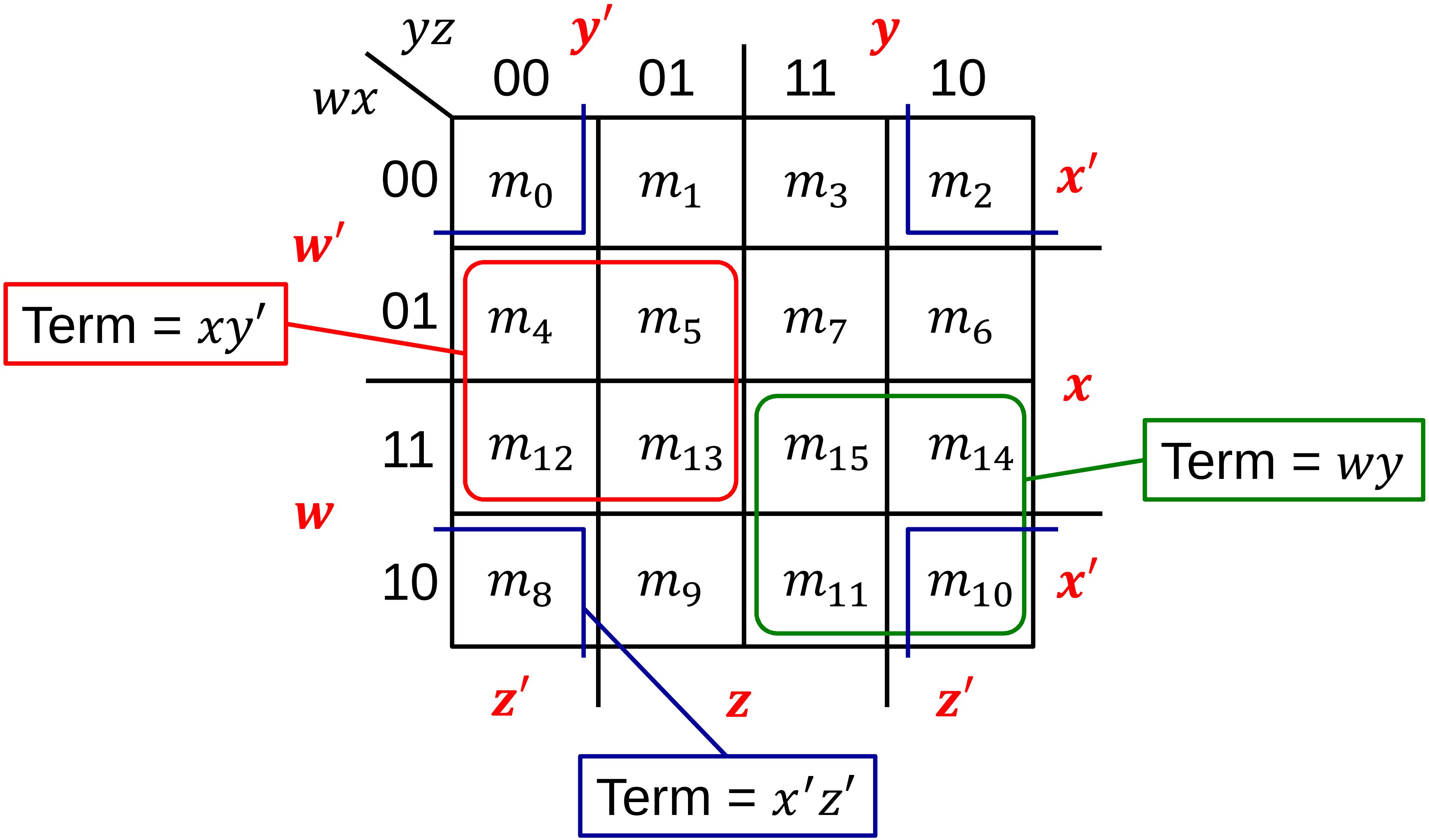
Combining Squares on a 4-Variable K-Map

- On a 4-variable K-Map:
  - One square represents a minterm with 4 variables
  - Two adjacent squares represent a term with 3 variables
  - Four adjacent squares represent a term with 2 variables
  - Eight adjacent squares represent a term with 1 variable
  - Combining all 16 squares is the constant **'1'** (no variables)

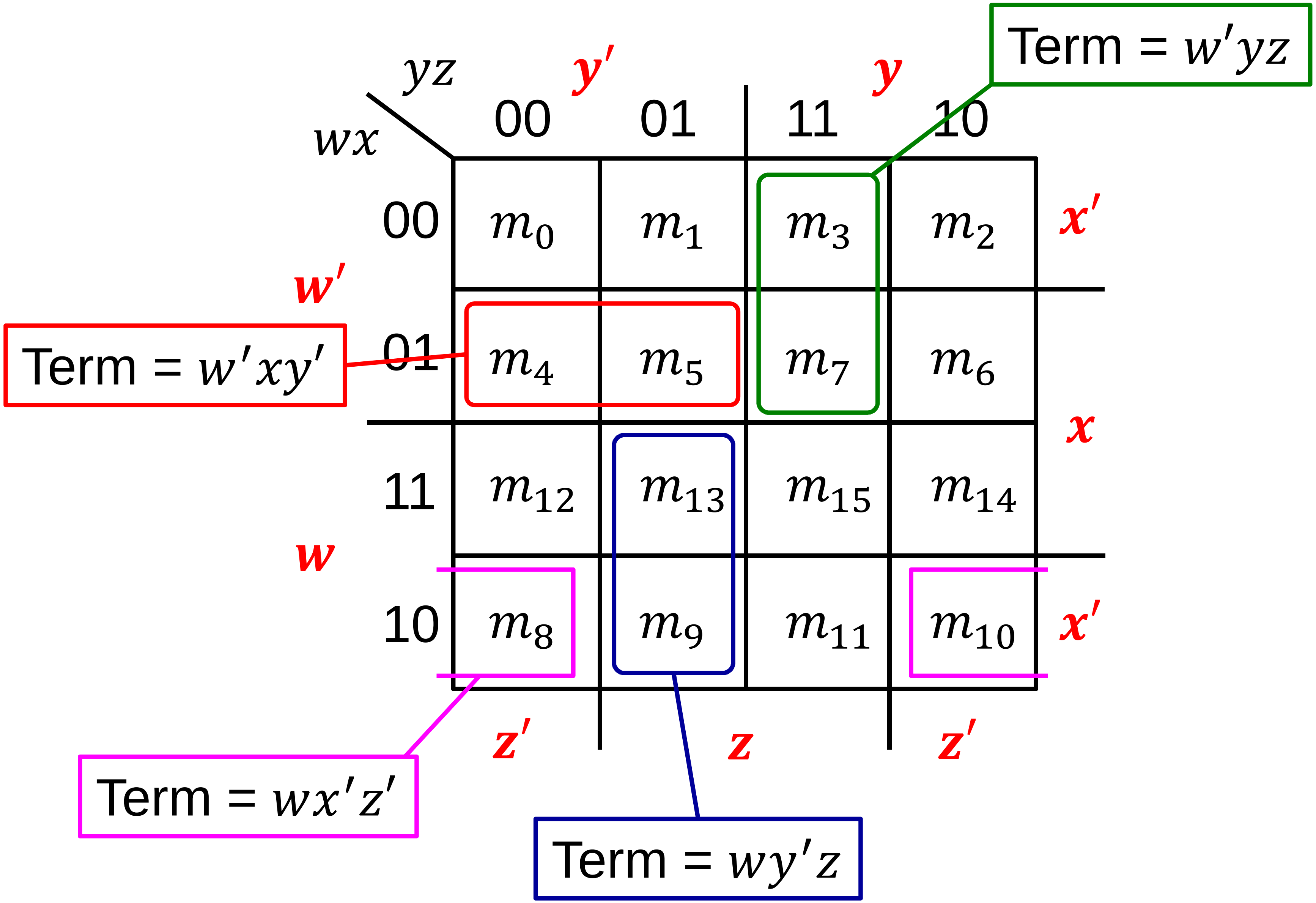
Combining Eight Squares



Combining Four Squares



Combining Two Squares



Simplifying a 4-Variable Function

Given $f(w, x, y, z) = \sum(0, 2, 4, 5, 6, 7, 8, 12)$

Draw the K-map for function f

Minimize f as sum-of-products

Solution:

$f = w'x + y'z' + w'z'$

$wx \backslash yz$	00	01	11	10
00	1	0	0	1
01	1	1	1	1
11	1	0	0	0
10	1	0	0	0

y' y

w' w

x' x x'

z' z z'

Term = $w'z'$

Term = $w'x$

Term = $y'z'$

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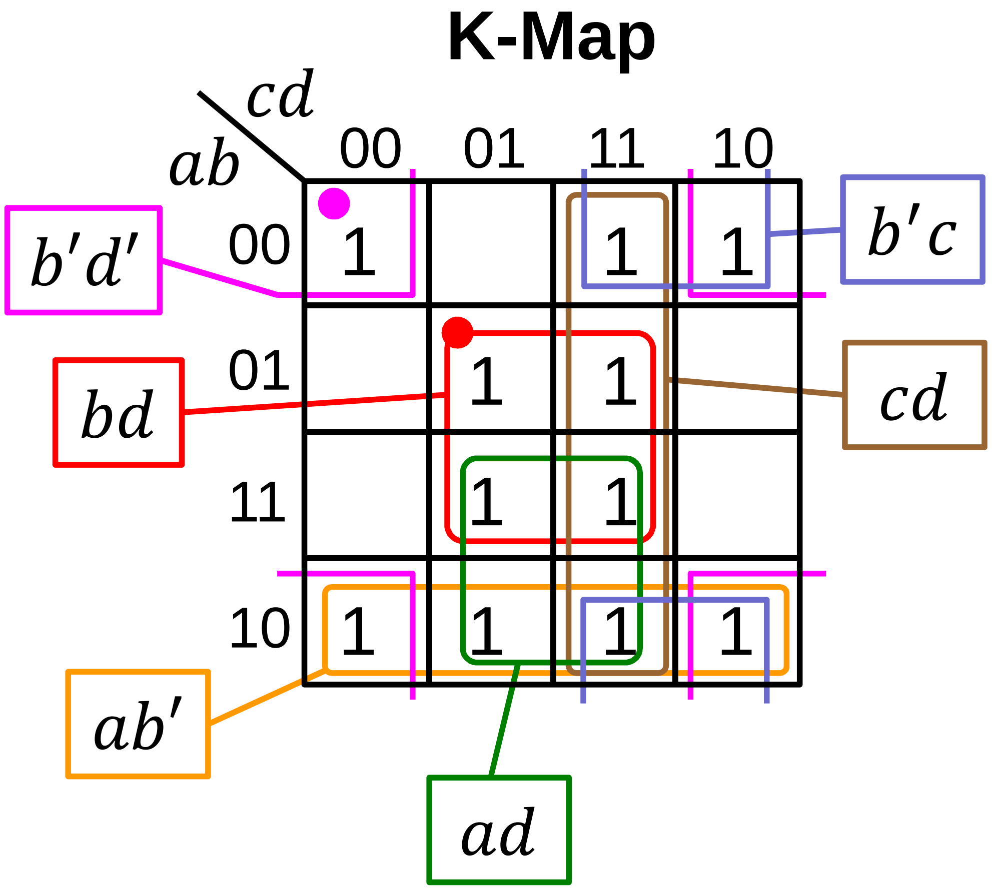
Prime Implicants

- **Prime Implicant:** a product term obtained by combining the **maximum number of adjacent squares** in the K-map
- The number of combined squares must be a **power of 2**
- **Essential Prime Implicant:** is a prime implicant that covers at least one minterm not covered by the other prime implicants
- The prime implicants and essential prime implicants can be determined by inspecting the K-map

Example of Prime Implicants

Find all the prime implicants and essential prime implicants for:

$$f(a,b,c,d) = \sum(0,2,3,5,7,8,9,10,11,13,15)$$



Six Prime Implicants

bd, $b'd'$, ab', ad, cd, $b'c$

Only Two Prime Implicants are essential

bd and $b'd'$

Simplification Procedure Using the K-Map

1. Find all the essential prime implicants
   - ✧ Covering maximum number (power of 2) of 1's in the K-map
   - ✧ Mark the minterm(s) that make the prime implicants essential
2. Add prime implicants to cover the function
   - ✧ Choose a minimal subset of prime implicants that cover all remaining 1's
   - ✧ Make sure to cover all 1's not covered by the essential prime implicants
   - ✧ Minimize the overlap among the additional prime implicants

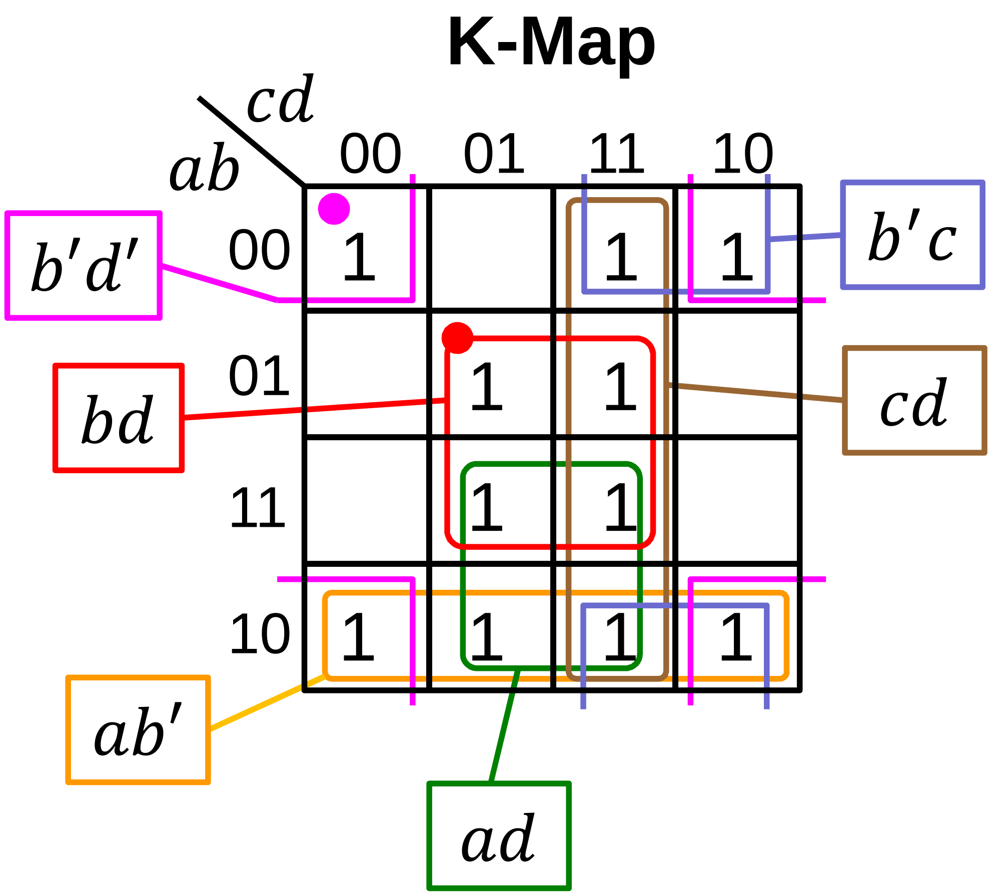
❖ Sometimes, a function has multiple simplified expressions

- ✧ You may be asked to list all the simplified sum-of-product expressions

Obtaining All Minimal SOP Expressions

Consider again: $f(a, b, c, d) = \sum(0, 2, 3, 5, 7, 8, 9, 10, 11, 13, 15)$

Obtain all minimal sum-of-products (SOP) expressions



Two essential Prime Implicants: bd and $b'd'$

Four possible solutions:

$f = bd + b'd' + cd + ad$

$f = bd + b'd' + cd + ab'$

$f = bd + b'd' + b'c + ab'$

$f = bd + b'd' + b'c + ad$

Product-of-Sums (POS) Simplification

- All previous examples were expressed in Sum-of-Products form
- With a minor modification, the Product-of-Sums can be obtained
- Example: $f(a, b, c, d) = \sum(1, 2, 3, 9, 10, 11, 13, 14, 15)$

K-Map of f

$ab \backslash cd$	00	01	11	10
00		1	1	1
01				
11		1	1	1
10		1	1	1

$f = ad + ac + b'd + b'c$

Minimal Sum-of-Products = 8 literals

All prime implicants are essential

K-Map of f'

$ab \backslash cd$	00	01	11	10
00	1			
01	1	1	1	1
11	1			
10	1			

$f' = c'd' + a'b$

$f = (c + d)(a + b') \rightarrow$ Minimal Product-of-Sums = 4 literals

Product-of-Sums Simplification Procedure

1. Draw the K-map for the function f
   - Obtain a minimal Sum-of-Products (SOP) expression for f
2. Draw the K-map for f', replacing the 0's of f with 1's in f'
3. Obtain a minimal Sum-of-Products (SOP) expression for f'
4. Use DeMorgan's theorem to obtain $f = (f')'$
   - The result is a minimal Product-of-Sums (POS) expression for f
5. Compare the cost of the minimal SOP and POS expressions
   - Count the number of literals to find which expression is minimal

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Don't Cares

- Sometimes, a function table may contain entries for which:
  - The input values of the variables will never occur, or
  - The output value of the function is never used
- In this case, the output value of the function is not defined
- The output value of the function is called a **don't care**
- A don't care is an **X** value that appears in the function table
- The **X** value can be later chosen to be **0 or 1**
  - To minimize the function implementation

Example of a Function with Don't Cares

- Consider a function f defined over BCD inputs
- The function input is a BCD digit from 0 to 9
- The function output is 0 if the BCD input is 0 to 4
- The function output is 1 if the BCD input is 5 to 9
- The function output is X (don't care) if the input is 10 to 15 (not BCD)
- $f = \underbrace{\sum_m (5, 6, 7, 8, 9)}_{\text{Minterms}} + \underbrace{\sum_d (10, 11, 12, 13, 14, 15)}_{\text{Don't Cares}}$

Truth Table

a b c d	f
0 0 0 0	0
0 0 0 1	0
0 0 1 0	0
0 0 1 1	0
0 1 0 0	0
0 1 0 1	1
0 1 1 0	1
0 1 1 1	1
1 0 0 0	1
1 0 0 1	1
1 0 1 0	X
1 0 1 1	X
1 1 0 0	X
1 1 0 1	X
1 1 1 0	X
1 1 1 1	X

Minimizing Functions with Don't Cares

Consider: $f = \sum_m (5,6,7,8,9) + \sum_d (10,11,12,13,14,15)$

If the don't cares were treated as 0's we get:

$f = a'bd + a'bc + ab'c'$ (9 literals)

If the don't cares were treated as 1's we get:

$f = a + bd + bc$ (5 literals)

The don't care values can be selected to be either 0 or 1, to produce a minimal expression

K-Map of f

$ab \backslash cd$	00	01	11	10
00	0	0	0	0
01	0	1	1	1
11	X	X	X	X
10	1	1	X	X

Simplification Procedure with Don't Cares

1. Find all the essential prime implicants
   - Covering maximum number (power of 2) of 1's and X's (don't cares)
   - Mark the 1's that make the prime implicants essential
2. Add prime implicants to cover the function
   - Choose a minimal subset of prime implicants that cover all remaining 1's
   - Make sure to cover all 1's not covered by the essential prime implicants
   - Minimize the overlap among the additional prime implicants
   - You need not cover all the don't cares (some can remain uncovered)

- Sometimes, a function has multiple simplified expressions

Minimizing Functions with Don't Cares (2)

Simplify: $g = \sum_m(1, 3, 7, 11, 15) + \sum_d(0, 2, 5)$

Solution 1: $g = cd + a'b'$ (4 literals)

Solution 2: $g = cd + a'd$ (4 literals)

Prime Implicant cd is essential

K-Map of g

$ab \backslash cd$	00	01	11	10
00	X	1	1	X
01	0	X	1	0
11	0	0	1	0
10	0	0	1	0

K-Map of g

$ab \backslash cd$	00	01	11	10
00	X	1	1	X
01	0	X	1	0
11	0	0	1	0
10	0	0	1	0

Not all don't cares need be covered

Minimal Product-of-Sums with Don't Cares

Simplify: $g = \sum_m(1,3,7,11,15) + \sum_d(0,2,5)$

Obtain a product-of-sums minimal expression

Solution: $g' = \sum_m(4,6,8,9,10,12,13,14) + \sum_d(0,2,5)$

Minimal $g' = d' + ac'$ (3 literals)

Minimal product-of-sums:

$g = d(a' + c)$ (3 literals)

The minimal sum-of-products expression for g had 4 literals

K-Map of g'

$ab \backslash cd$	00	01	11	10
00	X	0	0	X
01	1	X	0	1
11	1	1	0	1
10	1	1	0	1

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Five-Variable Karnaugh Map

- Consists of $2^5 = 32$ squares, numbered 0 to 31
  - Remember the numbering of squares in the K-map
- Can be visualized as two layers of 16 squares each
- Top layer contains the squares of the first 16 minterms ($a = 0$)
- Bottom layer contains the squares of the last 16 minterms ($a = 1$)

$a = 0$

$bc \backslash de$	00	01	11	10
00	m_0	m_1	m_3	m_2
01	m_4	m_5	m_7	m_6
11	m_{12}	m_{13}	m_{15}	m_{14}
10	m_8	m_9	m_{11}	m_{10}

$a = 1$

$bc \backslash de$	00	01	11	10
00	m_{16}	m_{17}	m_{19}	m_{18}
01	m_{20}	m_{21}	m_{23}	m_{22}
11	m_{28}	m_{29}	m_{31}	m_{30}
10	m_{24}	m_{25}	m_{27}	m_{26}

Each square is adjacent to **5** other squares:
4 in the same layer and
1 in the other layer:
m_0 is adjacent to m_{16}
m_1 is adjacent to m_{17}
m_4 is adjacent to m_{20} …

Example of a Five-Variable K-Map

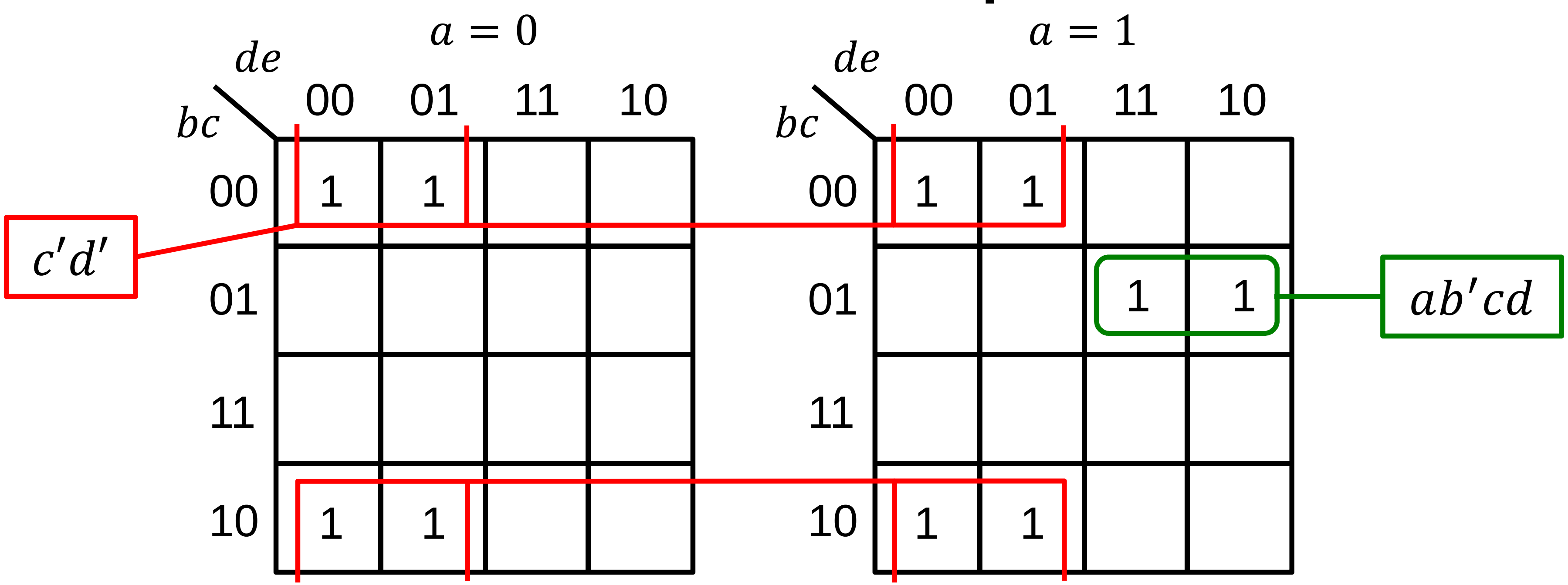
Given: $f(a,b,c,d,e) = \sum(0,1,8,9,16,17,22,23,24,25)$

Draw the 5-Variable K-Map

Obtain a minimal Sum-of-Products expression for f

Solution: $f \ = \ c'd' + ab'cd$ (6 literals)

5-Variable K-Map



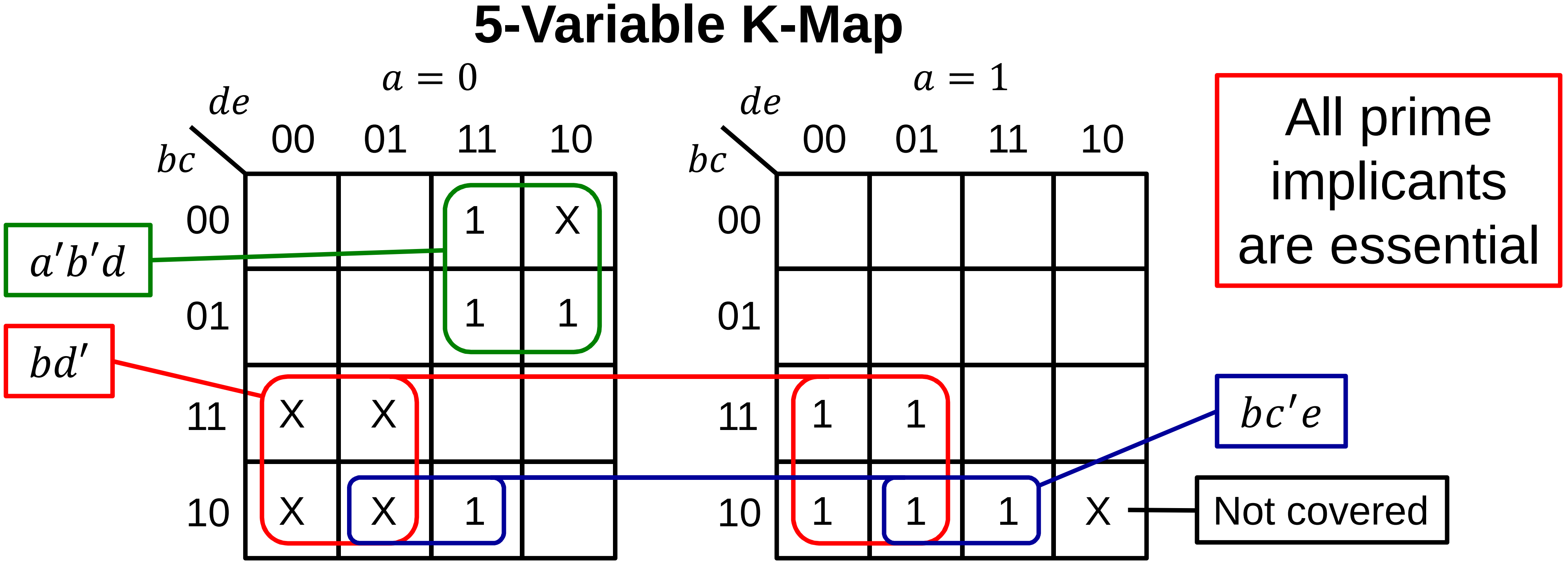
Five-Variable K-Map with Don't Cares

$g(a,b,c,d,e) = \sum_m(3,6,7,11,24,25,27,28,29) + \sum_d(2,8,9,12,13,26)$

Draw the 5-Variable K-Map

Obtain a minimal Sum-of-Products expression for g

Solution: $g = bd' + a'b'd + bc'e$ (8 literals)



Six-Variable Karnaugh Map

- Consists of $2^6 = 64$ squares, numbered 0 to 63
- Can be visualized as four layers of 16 squares each
  - Four layers: $ab = 00, 01, 11, 10$ (Notice that layer 11 comes before 10)
- Each square is adjacent to 6 other squares:
  - 4 squares in the same layer and 2 squares in the above and below layers

	$ab = 00$				$ab = 01$				$ab = 11$				$ab = 10$			
cd \ ef	00	01	11	10	00	01	11	10	00	01	11	10	00	01	11	10
00	m_0	m_1	m_3	m_2	m_{16}	m_{17}	m_{19}	m_{18}	m_{48}	m_{49}	m_{51}	m_{50}	m_{32}	m_{33}	m_{35}	m_{34}
01	m_4	m_5	m_7	m_6	m_{20}	m_{21}	m_{23}	m_{22}	m_{52}	m_{53}	m_{55}	m_{54}	m_{36}	m_{37}	m_{39}	m_{38}
11	m_{12}	m_{13}	m_{15}	m_{14}	m_{28}	m_{29}	m_{31}	m_{30}	m_{60}	m_{61}	m_{63}	m_{62}	m_{44}	m_{45}	m_{47}	m_{46}
10	m_8	m_9	m_{11}	m_{10}	m_{24}	m_{25}	m_{27}	m_{26}	m_{56}	m_{57}	m_{59}	m_{58}	m_{40}	m_{41}	m_{43}	m_{42}

Example of a Six-Variable K-Map

$h(a,b,c,d,e,f) = \sum(2, 10, 11, 18, 21, 23, 29, 31, 34, 41, 50, 53, 55, 61, 63)$

Draw the 6-Variable K-Map

Obtain a minimal Sum-of-Products expression for h

Solution: $h \;=\; c'd'ef' + b\,d\,f + a'b'c\,d'e + a\,b'\,c\,d'e'f$ (18 literals)

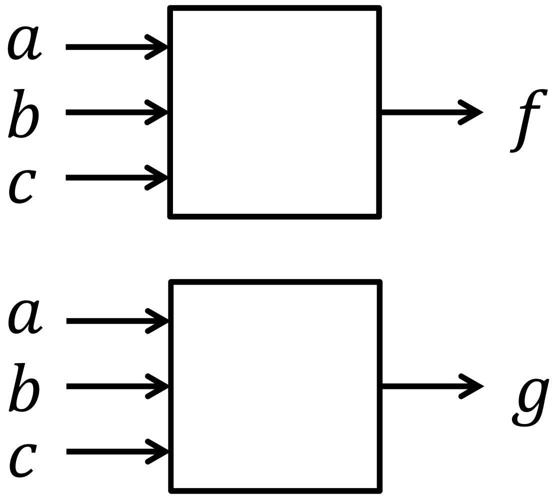
cd \ ef	ab = 00: 00	01	11	10	ab = 01: 00	01	11	10	ab = 11: 00	01	11	10	ab = 10: 00	01	11	10
00				1				1				1				1
01						1	1			1	1					
11						1	1			1	1					
10			1	1										1		

Next . . .

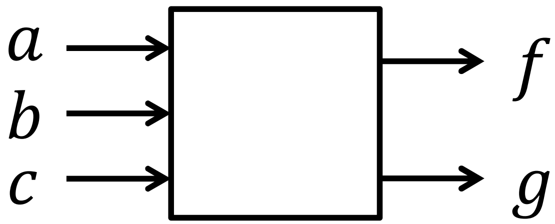
- ❖ Boolean Function Minimization
- ❖ The Karnaugh Map (K-Map)
- ❖ Two, Three, and Four-Variable K-Maps
- ❖ Prime and Essential Prime Implicants
- ❖ Minimal Sum-of-Products and Product-of-Sums
- ❖ Don't Cares
- ❖ Five and Six-Variable K-Maps
- ❖ Multiple Outputs

Multiple Outputs

- Suppose we have two functions: $f(a, b, c)$ and $g(a, b, c)$
- Same inputs: a, b, c, but two outputs: f and g
- We can minimize each function separately, or
- Minimize f and g as one circuit with two outputs
- The idea is to share terms (gates) among f and g



Two separate circuits



One circuit with Two Outputs

Multiple Outputs: Example 1

Given: $f(a,b,c) = \sum(0,2,6,7)$ and $g(a,b,c) = \sum(1,3,6,7)$

Minimize each function separately

Minimize both functions as one circuit

K-Map of f

$a \backslash bc$	00	01	11	10
0	1	0	0	1
1	0	0	1	1

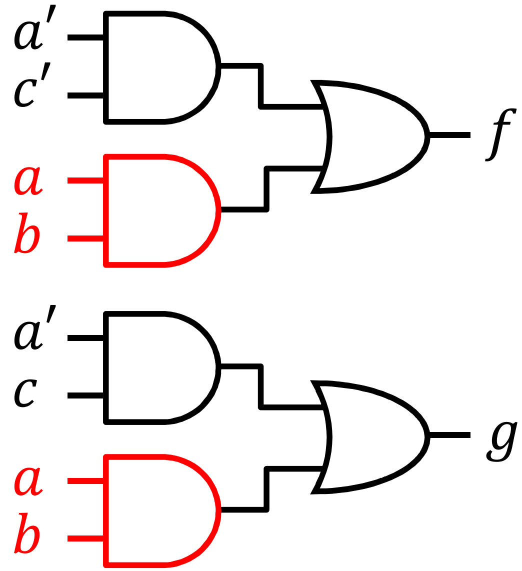
$f = a'c' + ab$

K-Map of g

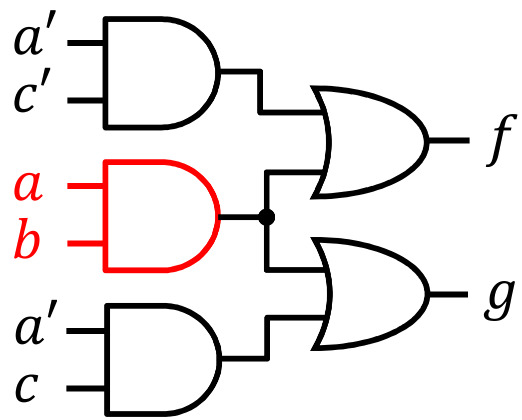
$a \backslash bc$	00	01	11	10
0	0	1	1	0
1	0	0	1	1

$g = a'c + ab$

Common Term = ab



One circuit per function



One circuit with two Outputs

Multiple Outputs: Example 2

$f(a,b,c,d) = \sum(3,5,7,10,11,14,15)$, $g(a,b,c,d) = \sum(1,3,5,7,10,14)$

Draw the K-map and write minimal SOP expressions of f and g

$f = a'bd + ac + cd$ $\quad\quad g = a'd + acd'$

Extract the common terms of f and g

K-Map of f

$ab \backslash cd$	00	01	11	10
00			1	
01		1	1	
11			1	1
10			1	1

K-Map of g

$ab \backslash cd$	00	01	11	10
00		1	1	
01		1	1	
11				1
10				1

Common Terms

$T_1 = a'd$ and $T_2 = ac$

Minimal f and g

$f = T_1 b + T_2 + cd$

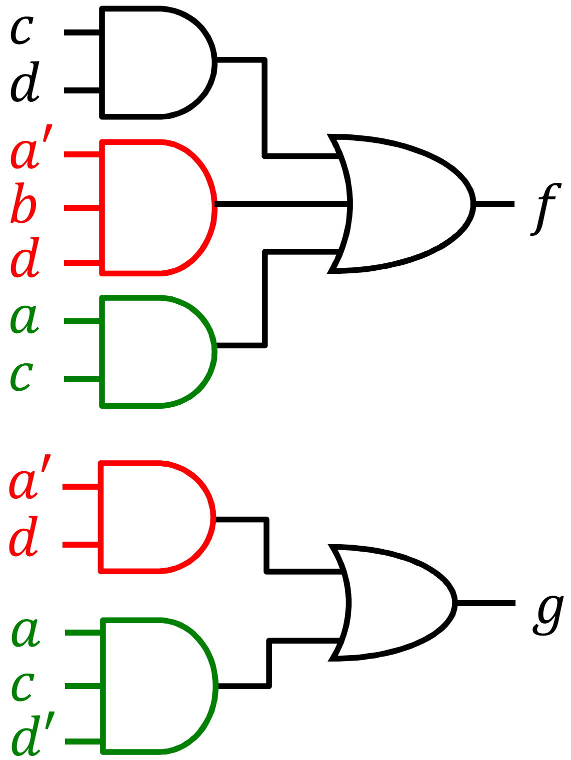
$g = T_1 + T_2 d'$

Common Terms → Shared Gates

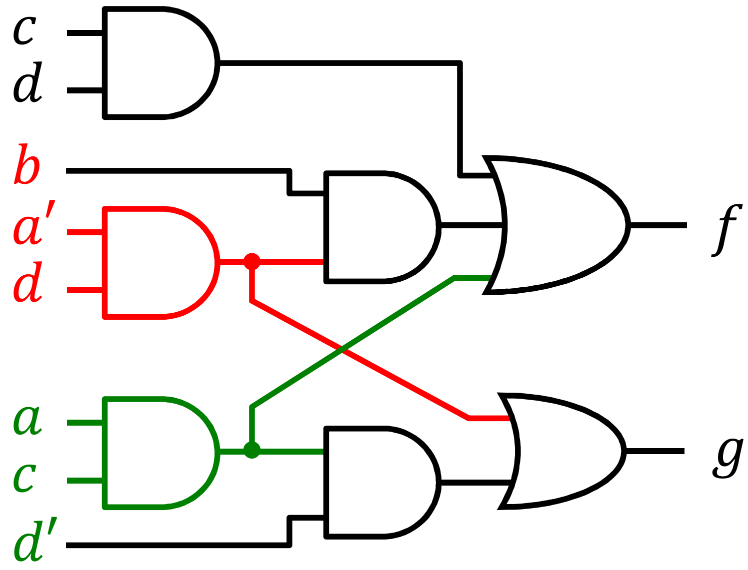
Minimal $f = a'bd + ac + cd$ Minimal $g = a'd + acd'$

Let $T_1 = a'd$ and $T_2 = ac$ (shared by f and g)

Minimal $f = T_1 b + T_2 + cd$, Minimal $g = T_1 + T_2 d'$



NO Shared Gates



One Circuit

Two Shared Gates